

# *Data Structures & Algorithms for Geometry*

## ⇒ Agenda:

- Robustness of calculations
- Prepare for final
- Assignment #4

# Fixed-point Representation

- ⇒ Select an  $n$ -bit data size, partition  $k$ -bits for the integral part and  $(n-k)$ -bits for the fractional part
  - Numbers are evenly distributed
    - The difference between all representable numbers is always  $1/(2^{n-k})$ .
  - Color is usually done this way: 0-bits for the integral part, 8-bits for the fractional part
  - Also the way *all* real-time graphics math is done on processors without floating-point units
    - Pre-486, MIPS CPU in Playstation, Pre-68040, many “embedded” CPUs

# Floating-point Representation

⇒ Single precision:

- 1-bit for sign
- 8-bits for exponent
- 23-bits for fractional part

⇒ If  $E \neq 0$ , value =  $(-1^S \times 1.F \times 2^{E-127})$

⇒ If  $E = 0$ , value =  $(-1^S \times 0.F \times 2^{-126})$

- These are *denormalized* (denorm for short)

# *Magic Values*

- ⇒ IEEE-754 defines some magic values:
  - $(E = 0, F = 0, S = 0) \rightarrow 0$
  - $(E = 0, F = 0, S = 1) \rightarrow -0$
  - $(E = 255, F = 0, S = 0) \rightarrow +\infty$
  - $(E = 255, F = 0, S = 1) \rightarrow -\infty$
  - $(E = 255, F \neq 0) \rightarrow$  Not a number (a.k.a., NaN)
    - More on NaN and Inf in a few slides

# *Distribution of Values*

⇒ As the exponent increases, the real difference between two values increases

- $(E = 200, F = 1) - (E = 200, F = 0) \rightarrow \sim 4 \times 10^{56}$

- $(E = 1, F = 1) - (E = 1, F = 0) \rightarrow \sim 1.4 \times 10^{-45}$

⇒ Usually, this is okay.

- This matches the notion of *significant digits*

# *NaN and Inf*

⇒  $\pm\infty$  are used to represent overflow cases

- $1.0e20 / 1.0e-20 = +\text{Inf}$
- $-1.0e20 / 1.0e-20 = -\text{Inf}$
- $1.0 / 0.0 = +\text{Inf}$

⇒ Not-a-number (NaN) is used to represent incalculable cases

- $0 / 0 = \text{NaN}$
- $\sqrt{-1} = \text{NaN}$
- $\text{Inf} - \text{Inf} = \text{NaN}$

# NaN Quirks

- ⇒ All comparisons involving NaN are false.
  - Except  $\text{NaN} \neq \text{NaN}$ , which is true.
- ⇒ This means the following code segments will produce *different* results if  $x$  is NaN!

```
if (X < value) { a() } else { b() }
```

```
...
```

```
if (X >= value) { b() } else { a() }
```

# NaN Quirks (cont.)

⇒ NaN can *help* by avoiding the need for divide by zero tests

- This code from the textbook works even if  $\text{Dot}(p.n, ab)$  is zero

```
int IntersectSegmentPlane(Point a, Point b, Plane p,
    float &t, Point &q)
{
    Vector ab = b - a;
    t = (p.d - Dot(p.n, a)) / Dot(p.n, ab);
    if (t >= 0.0f && t <= 1.0f) {
        q = a + t * ab;
        return 1;
    }
    return 0;
}
```



# References

Hecker, Chris. 1996. Let's Get to the (Floating) Point. Game Developer Magazine. (Feb. / Mar. 1996).

[http://chrishecker.com/Miscellaneous\\_Technical\\_Articles#Floating\\_Point](http://chrishecker.com/Miscellaneous_Technical_Articles#Floating_Point)

Goldberg, D. 1991. What every computer scientist should know about floating-point arithmetic. ACM Computing Surveys. 23, 1 (Mar. 1991), 5-48.

[http://docs.sun.com/source/806-3568/ncg\\_goldberg.html](http://docs.sun.com/source/806-3568/ncg_goldberg.html)

<http://en.wikipedia.org/wiki/IEEE-754>

# *Conversion and Representation Errors*

- ⇒ *Lots* of interesting, real numbers cannot be exactly represented
  - The more significant digits involved, the more *inexact* the representation will be
  - $\sqrt{2}$  get rounded to 1.41421356237309504880
    - This is an irrational number with infinite significant digits

# *Overflow and Underflow Errors*

- ⇒ If the result is too large (or too small) the result will overflow (or underflow)
  - Multiply two very large (or small) numbers
  - Divide a small number by a large number (or vice versa)
  - Overflows will result in  $\pm\text{Inf}$
  - Underflows will result in  $\pm 0$

# Round-off Errors

⇒ Result of an operation has more significant digits than can be represented

- $X = (E = 3, F = 1.413351774) = 11.306814194$

- $Y = (E = 7, F = 1.933333278) = 247.466659546$

- Results of  $X * Y$ :

- True result:

- **2798.05953862742171622812747955322265625**

- Computer result:  $(E = 10, F = 2.732480049) = (E = 11, F = 1.366240025) = \mathbf{2798.0595703125}$

# Digit-cancellation Errors

## ⇒ Related to round-off errors

- Subtracting nearly equal values
- Adding or subtracting a large value and a small value
  - $1e20 + 1e-20 = 1e20$

## ⇒ Floating point arithmetic is *not* associative!

- $(9876543.0 + -9876547.0) + 3.45 = -0.5499999...$
- $9876543.0 + (-9876547.0 + 3.45) = -1.0$

# *Input Errors*

- ⇒ The source data may have errors
  - Inexact measurement from a physical device
  - Errors from previous calculations
  - etc.

# *References*

[http://www.mpi-inf.mpg.de/~kettner/pub/nonrobust\\_ecgtr\\_04\\_a.html](http://www.mpi-inf.mpg.de/~kettner/pub/nonrobust_ecgtr_04_a.html)

[http://en.wikipedia.org/wiki/Floating\\_point](http://en.wikipedia.org/wiki/Floating_point)

# *Robust Comparisons*

- ⇒ Obviously, direct comparison for equality are just plain wrong



# *Robust Comparisons*

- ⇒ Obviously, direct comparison for equality are just plain wrong
- ⇒ First improvement: compare absolute difference to some small value,  $\varepsilon$

```
if (fabs(x - y) < epsilon) { ... }
```

- Called *absolute tolerance*
- Picking  $\varepsilon$  that works for a range of values is difficult or impossible
  - `sqrt(FLT_EPSILON)` is usually a good choice

# Robust Comparisons (cont.)

⇒ Second improvement: compare abs ratio to 1.0

- Called *relative tolerance*

```
if (fabs((x / y) - 1.0) <= epsilon)
```

- Assumes  $|x| < |y|$

```
if (fabs((x - y) / y) <= epsilon)
```

- First re-write

```
if (fabs(x - y) <= epsilon * fabs(y))
```

- Eliminate division

```
if (fabs(x - y) <= epsilon * max(fabs(x),  
    fabs(y)))
```

- Eliminate  $|x| < |y|$  assumption

# *Robust Comparisons (cont.)*

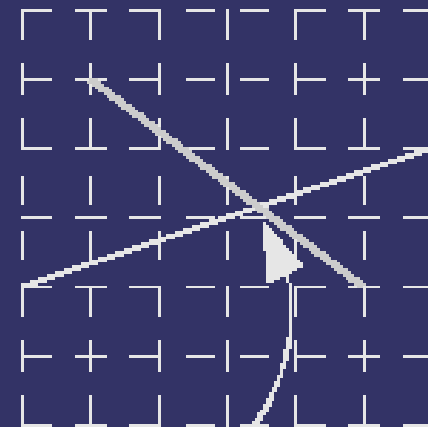
⇒ Tolerances do have one ugly side-effect

## *Robust Comparisons (cont.)*

- ⇒ Tolerances do have one ugly side-effect
  - $A = B$  and  $B = C$  does *not* imply  $A = C$

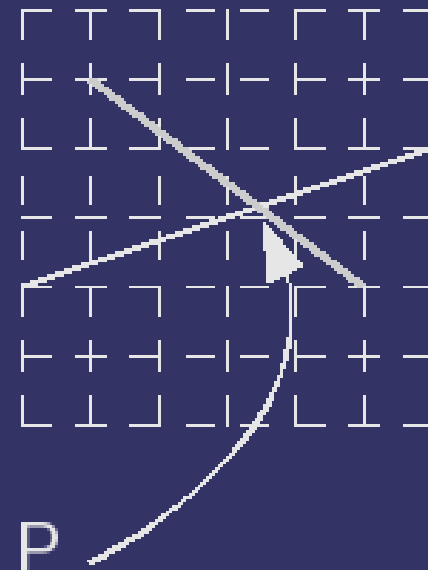
# *Thick Planes Revisited*

- ⇒ Not all numbers can be exactly represented in floating-point, so the true intersection point of a line and a plane may not be representable
  - Other object-object intersections similarly affected
  - Clipping the line to the plane effectively moves part of the line



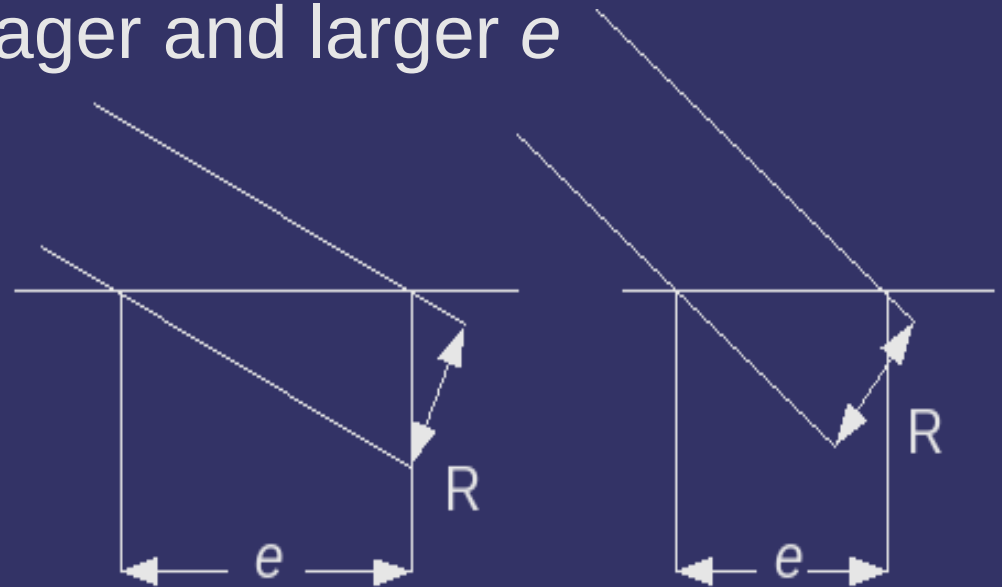
# Thick Planes Revisited

- ⇒ Let  $e$  be the maximum error in the calculated  $P$ , then the thick plane has radius  $r$ ,  $r > e$ 
  - How do we select an appropriate  $e$ ?



# Thick Planes Revisited

- ⇒ Let  $e$  be the maximum error in the calculated  $P$ , then the thick plane has radius  $r$ ,  $r > e$ 
  - How do we select an appropriate  $e$ ?
  - As the line and plane become more parallel, a selection of  $R$  leads to larger and larger  $e$
  - The selection of  $R$  limits the size of line segments or polygons that we can track
    - Pick  $R$  based on the size of the smallest line or polygon



# *Next week...*

## ⇒ Final:

- **Tuesday, December 11<sup>th</sup> at 7:45PM.**
- **DO NOT BE LATE!**



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